

# IMPROVED INTEGER PROGRAMMING-BASED NEIGHBORHOOD SEARCH FOR LTL LOAD PLAN DESIGN

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# WHAT TO REMEMBER

- ① Load planning for LTL trucking, a problem of *service network design*, has grown in complexity
- ② New tools needed for load planning
  - ▶ Faster transit times
  - ▶ Service requirements
- ③ Our approach
  - ▶ Detailed time-expanded network model
  - ▶ Very large path-selection IP with intree constraints
  - ▶ IP-based neighborhood search solution approach
- ④ Results
  - ▶ *Intree* neighborhood search: 3-5% \$ savings
  - ▶ *New or old tree* neighborhood search: 35-40% \$ savings

# WHAT DOES AN LTL CARRIER DO?

- Transports shipments from multiple origins to destinations
  - ▶ Individual shipments are small, do not fill trailer (“less-than-truckload”)
  - ▶ Each shipment has service requirement (# of business days)

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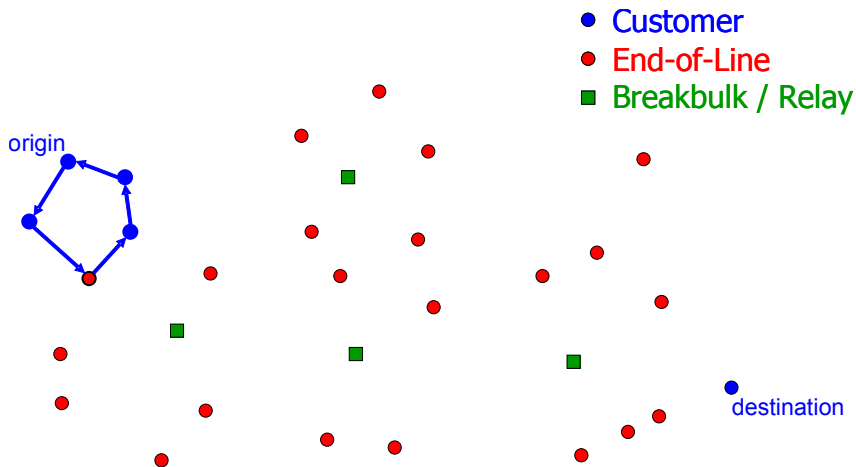
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- Consolidates shipments to reduce costs
  - ▶ Two terminal types
    - ★ End-of-lines (EOL - “Spoke”)
    - ★ Breakbulks (BB - “Hub”)
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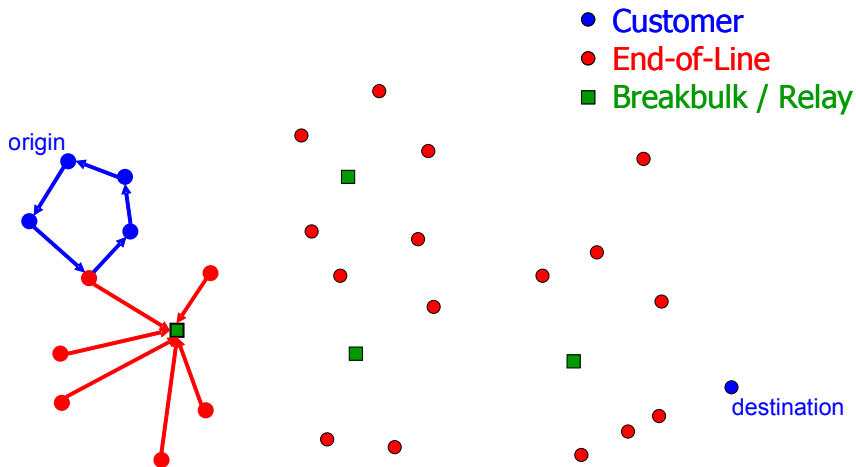
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    - ★ Handling cost
    - ★ Requires time, from 30 min to a few hours
- High-volume operations
  - ▶ Each week, a large US carrier
    - ★ Moves hundreds of millions lbs freight
    - ★ Hauls trailers millions of miles
    - ★ Spend millions of dollars on linehaul operation

# LTL OPERATIONS

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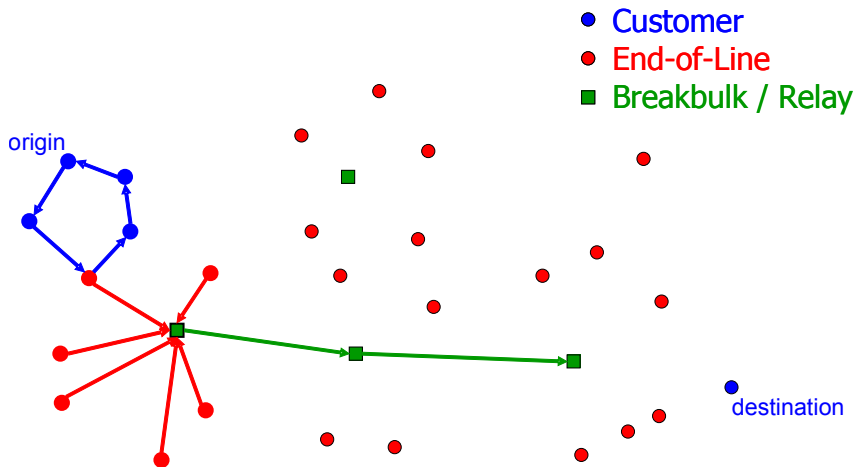


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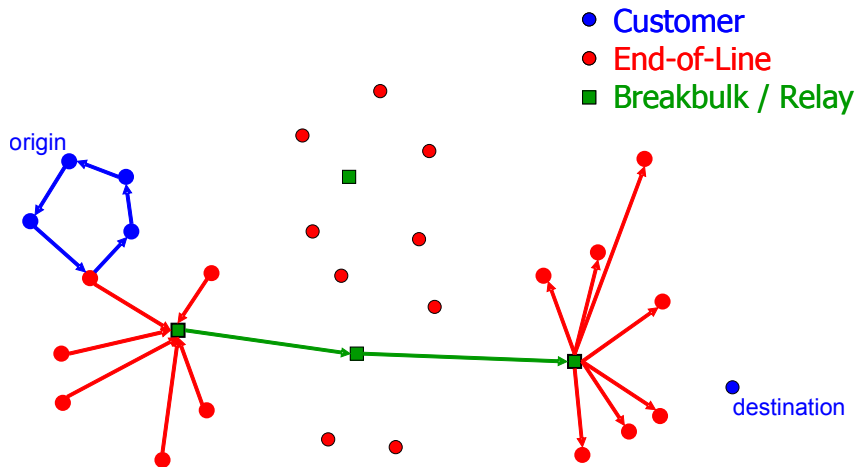




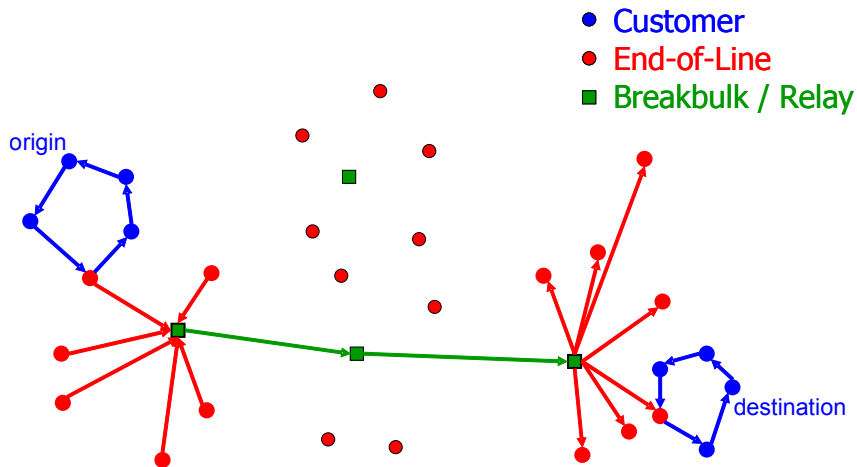
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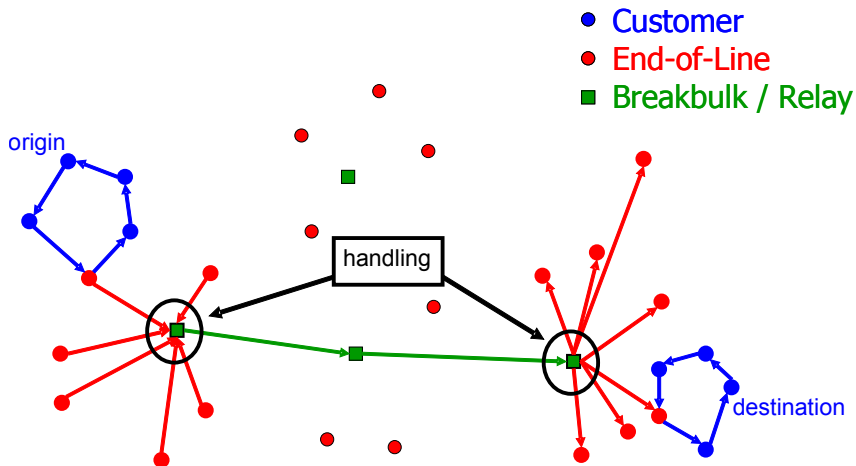
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# LTL LINEHAUL SERVICE NETWORK DESIGN

## LOAD PLANNING

Given:

- Terminal locations
- Terminal types
- Transportation, handling cost structure

Determine:

- Unique freight transfer path for each (origin, destination) terminal pair

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## LOOSELY-SCHEDULED OPERATIONS

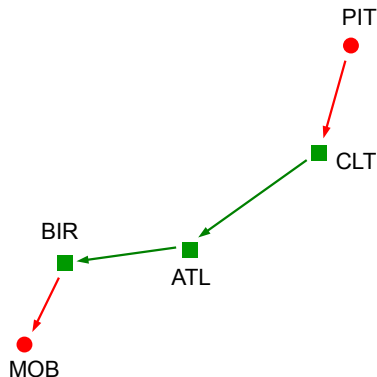
Given a load plan, freight arriving at terminals loaded into trailers, then dispatched using scheduled drivers, on-call “extra board” drivers, and third-party outsourcing (rail)

# LTL LOAD PLANNING

## DEFINITION

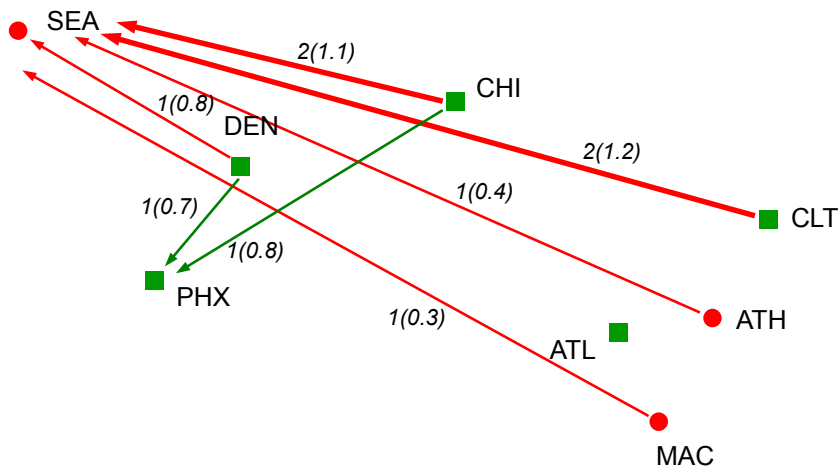
### Freight transfer paths

- A freight path is a sequence of direct trailer moves (*directs*)
- Freight assigned to direct  $(A, D)$  is loaded into a trailer at  $A$  and unloaded for transfer at  $D$



# OVERLAPPING FREIGHT PATHS AND CONSOLIDATION

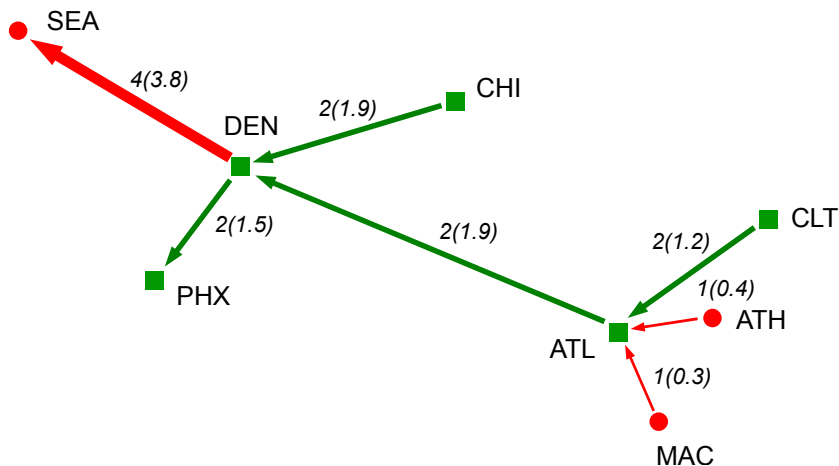
Simple illustration





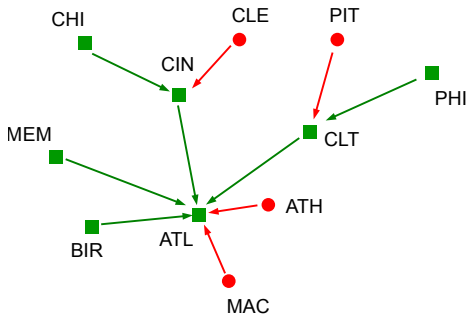
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Simple illustration



## LOAD PLAN IN-TREE STRUCTURE

- Set of directs forming freight paths to destination terminal  $d$  is directed intree
- Triples  $(i, d, j)$  for all  $i, d$  define load plan
  - ▶ All originating *and* transfer freight at  $i$  bound for final destination  $d$  is loaded into a trailer bound for next destination  $j$



# MODELS SUPPORTING LTL LOAD PLANNING

## Deterministic tactical planning models

- Flat (static) network, sequential empties
  - ▶ Powell and Sheffi (1983, 1989), Powell (1986)
  - ▶ Powell and Koskosidis (1992)
- Dynamic network, sequential empties, nointree
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- **Detailed dynamic network, integrated empties**
  - ▶ Erera, Hewitt, Savelsbergh, Zhang (under review)
  - ▶ Erera, Lindsey, Savelsbergh (working paper)

# OVERVIEW OF OUR APPROACH

One week wrapped dynamic network with overnight detail

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## WHY?

- More accurately model consolidation timing
- Input O-D freight volume to vary by day-of-week
- Properly represent scheduled outsourced transportation (rail)
- Model weekend operations

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Path-based freight transfer path selection IP



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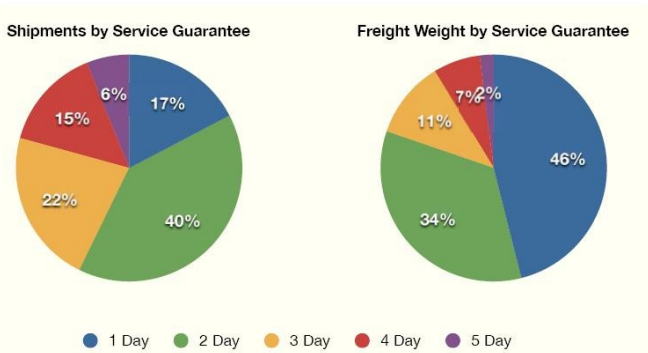
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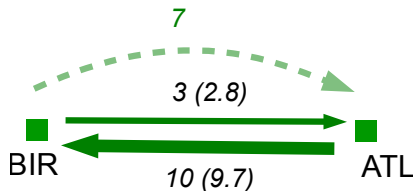
- Input freight imbalance creates implicit empty movement “demand”
- Creates opportunities for low-cost transfer of backhaul freight

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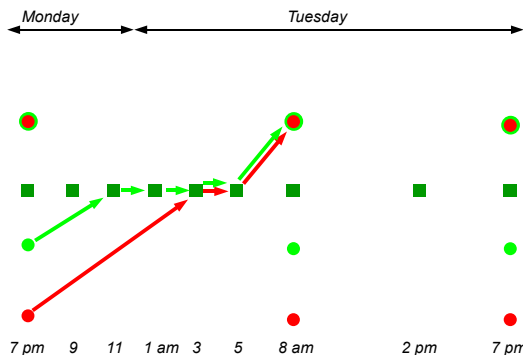
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# DYNAMIC NETWORK AND FREIGHT PATHS

## DYNAMIC NETWORK

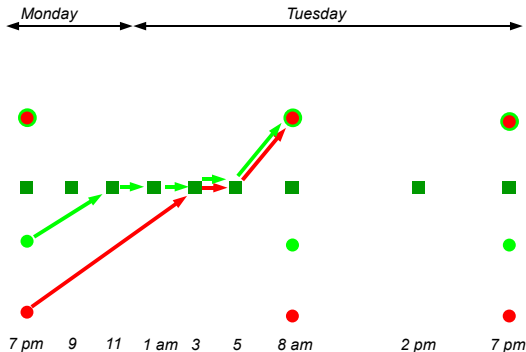
- Node: Terminal and time
- Arc: Freight direct or holding, or empty trailer move or holding
- Commodity: Freight demand from  $(o, t_1)$  to  $(d, t_2)$  each day



# DYNAMIC NETWORK AND FREIGHT PATHS

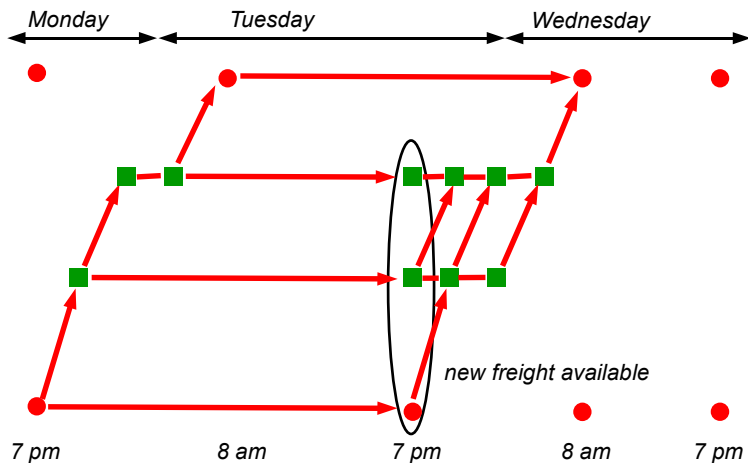
## FEASIBLE FREIGHT PATH GENERATION

- $m$  minimum cost flat network paths
- timed variants of each flat network path



# TIMED FREIGHT PATHS

*An two-day commodity originating Monday with slack time*





# OPTIMIZATION MODEL

Linear multi-commodity flow IP with side constraints

## OBJECTIVE FUNCTION

Minimize sum of:

- Transportation costs (linear in integer trailer variables)
- Holding costs (linear in integer path variables)

## CONSTRAINTS

- Select one path for each commodity
- Ensure that a single outbound direct is selected from each terminal  $i$  for freight destined to  $d$  [or, one direct for each  $i, wd(i)$ ]
- Only allow selection of paths if all component directs are selected
- Count required trailers for each direct
- Ensure trailer count balance at all nodes (empty repositioning)

## LOAD PLAN IP

$$\min \quad \sum_{a \in A} c_a \tau_a + \sum_{k \in K} \sum_{p \in P(k)} h_p q_k x_p^k$$

subject to

$$\sum_{p \in P(k)} x_p^k = 1 \quad \forall k \in K$$

$$\sum_{\ell \in \delta^+(u)} y_\ell^d \leq 1 \quad \forall u \in U, \forall d \in U$$

$$\sum_{p \in P(k): a \in p} x_p^k \leq y_{\ell(a)}^{d(k)} \quad \forall k \in K, \forall a \in A$$

$$\sum_{k \in K} \sum_{p \in P(k): a \in p} q_k x_p^k \leq \tau_a \quad \forall a \in A$$

$$\sum_{a \in \delta^+(v)} \tau_a - \sum_{a \in \delta^-(v)} \tau_a = 0 \quad \forall v \in V$$

# GIGANTIC REAL-WORLD INSTANCES

Saia, a large U.S. super-regional LTL carrier

- 93 EOLs
- 61 transfer terminals (BBs and smaller transfer points)
- Flat network directs  $|L| = 15,658$
- Time-space nodes  $|N| = 4,718$
- Time-space arcs  $|A| = 502,946$

Instance	# Commodities $ K $	# Paths $\sum_{k \in K}  P(k) $
Mar09-W1	36,135	2,978,300
Mar09-W2	36,218	2,979,323
Mar09-W3	36,599	2,981,454
Mar09-W4	36,783	2,982,673

# INTEGER PROGRAMMING BASED LOCAL SEARCH

## IP-BASED LS USING SINGLE INTREE NEIGHBORHOOD

**Require:** Current load plan solution  $(\bar{y}, \bar{x}, \bar{\tau})$

**while** search time has not exceeded time limit  $T$  **do**

    Choose a destination terminal  $d$

    Solve\* Single Intree IP for  $d$

**if** New solution has lower cost than current **then**

        Update  $(\bar{y}, \bar{x}, \bar{\tau})$

**end if**

**end while**

# SEARCHING SINGLE INTREE NEIGHBORHOOD

## SINGLE INTREE IP

Given single  $d$  and  $K(d)$ :

- Fix  $y = \bar{y}$ ,  $x = \bar{x}$  in Load Plan IP, except those involving  $d$  or  $k \in K(d)$
- Compute arc freight flows  $f_a$
- Create Load Plan IP with only variables for  $d$  and  $K(d)$
- Add trailer bounds  $\tau_a \geq \lceil f_a \rceil$

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## SOLVING SINGLE INTREE IP

- CPLEX 11 (or Gurobi)
- MIPemphasis = *integer feasibility*
- Gap 0.1%, short time limit  $T' \ll T$

# RESULTS AND IDEAS FOR IMPROVEMENT

## IP Based LS using Single Intree Neighborhood Results

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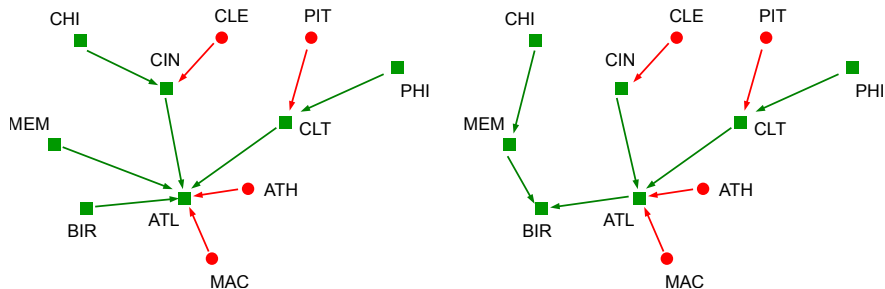
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### NEW IDEAS

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- Replace empty balance constraints with trailer flow bounds updated periodically within search

# MULTIPLE DESTINATION INTERACTIONS

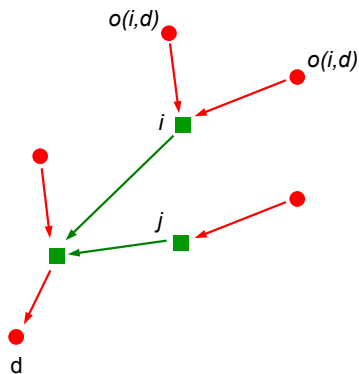


When changing freight paths into ATL, some paths into BIR should be simultaneously changed

# NEW OR OLD TREE NEIGHBORHOOD

## SETS OF INTERACTING O-D PAIRS

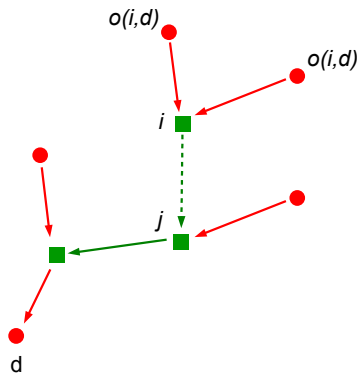
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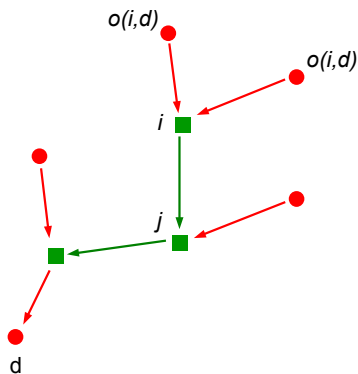
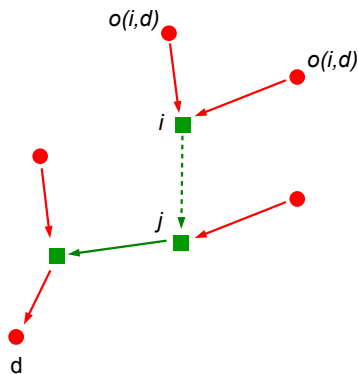




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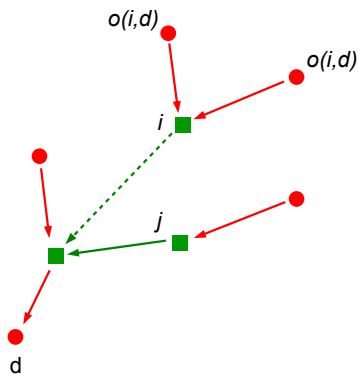
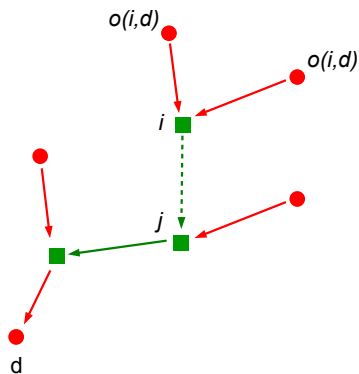
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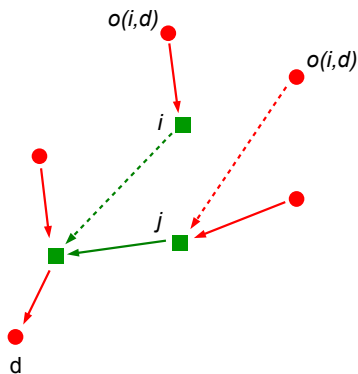
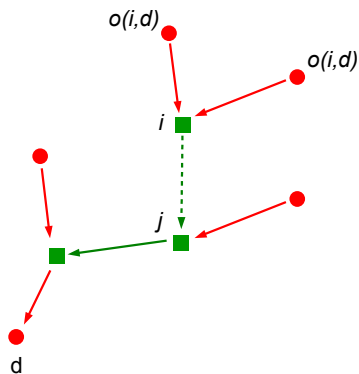
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subject to

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$$x_p^k \leq 1 - z_d \quad \forall d \in D', k \in K'(d), p \in P(k, \text{NewTree}(d))$$

$$x_p^k \leq z_d \quad \forall d \in D', k \in K'(d), p \in P(k, \text{OldTree}(d))$$

$$\sum_{k \in K'} \sum_{p \in P(k)} q_k x_p^k + f_a \leq \tau_a \quad \forall a \in A$$

$$x_p^k, z_d \text{ binary}$$

$$\tau_a \geq \text{Empties}_a \text{ and integer}$$

# EMPTY BALANCING OUTSIDE OF IP

Substantially reduce IP size within neighborhood search

## PERIODIC EMPTY BALANCE

- Every  $k$  neighborhood search iterations, solve empty balance MCNF on time-space network
- Let  $Empties_a$  be number of empty trailers moving on time-space arc  $a$

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    Choose whether to attract or reduce freight  
    Choose a direct  $(i, j)$   
    Solve\* New or Old Tree IP  
    **if** New solution has lower cost than current **then**  
        Update  $(\bar{y}, \bar{x}, \bar{\tau})$   
    **end if**  
    **if** iterations since last empty balance exceeds threshold **then**  
        Solve Empty Balance MCFN generating new  $Empties_a$   
    **end if**  
**end while**

# NEW RESULTS

## ① Search IP Speed

- ▶ *Intree*: 90s limit, 75% to optimality
- ▶ *New Or Old Tree*: 30s limit, 97% to optimality

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		RO	AO	Rand(0.25)
Mar09-W1	$\Delta Cost$	20.41%	17.88%	24.07%
Mar09-W2	$\Delta Cost$	19.28%	16.76%	22.45%
Mar09-W3	$\Delta Cost$	18.16%	17.38%	22.24%
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## ③ Cost Reduction from Base Plan: 6 hours

		RO	AO	Rand(0.25)
Mar09-W1	$\Delta Cost$	28.96%	24.47%	38.79%
Mar09-W2	$\Delta Cost$	29.14%	24.76%	38.60%
Mar09-W3	$\Delta Cost$	28.44%	24.80%	38.16%
Mar09-W4	$\Delta Cost$	28.47%	23.13%	38.48%

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  - ▶ More accurate cost estimator for load plan

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- ③ Daily Load Plan Adjuster
  - ▶ Multiple daily updates for load plans given freight pickups
  - ▶ GRASP

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